

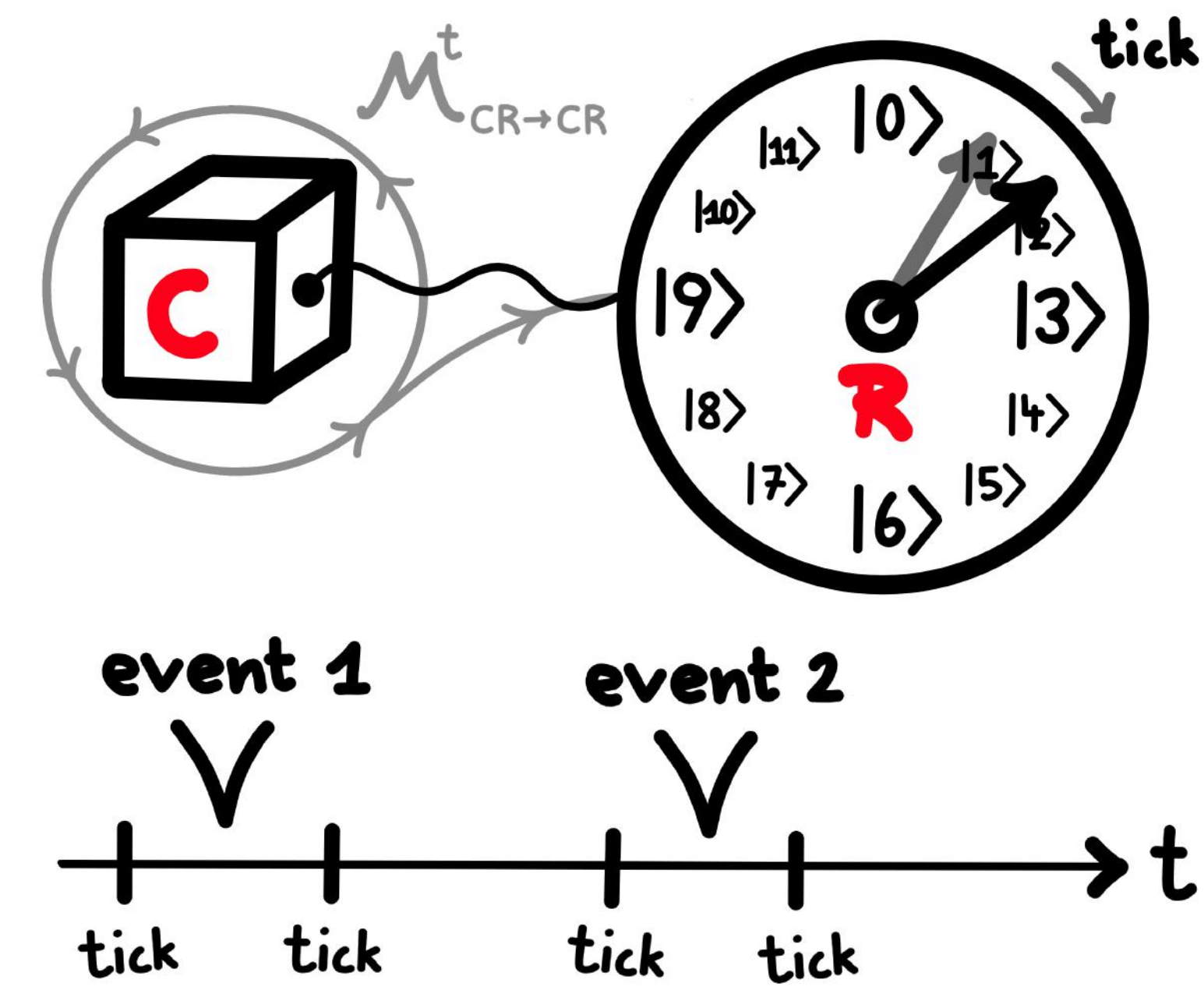
# Entropy production in ticking clocks: Fundamental limits of timekeeping

Julian Arnold, Mischa P. Woods

## Ticking clock model

A ticking clock is a timekeeping device that **autonomously** outputs time information in the form of individual ticks

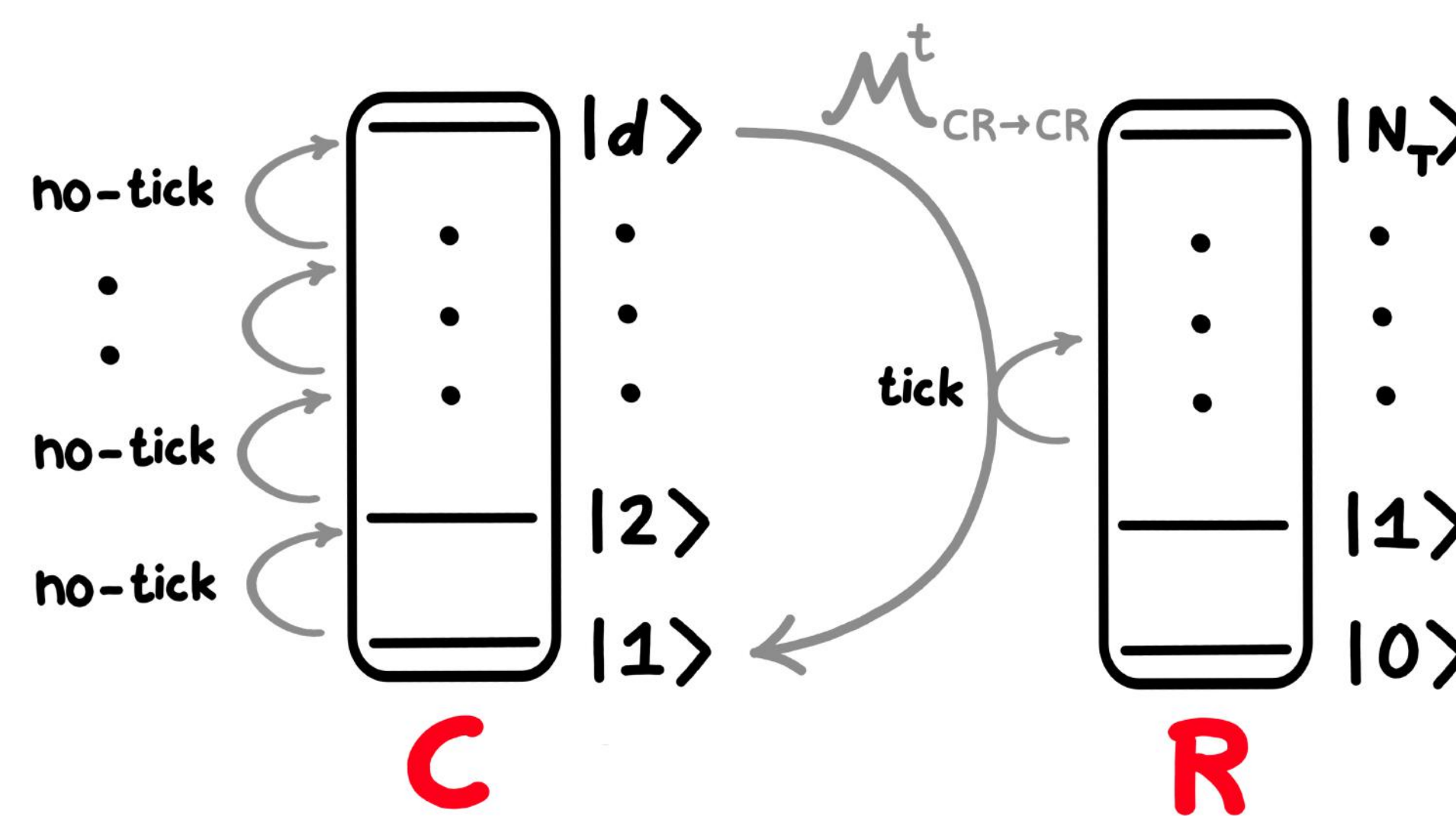
- modelled as a bipartite quantum system  $\rho_{CR}$  composed of clockwork **C** and register **R** living in  $\mathcal{H}_C \otimes \mathcal{H}_R$  with  $\dim(\mathcal{H}_C) = d$  and  $\dim(\mathcal{H}_R) = N_T + 1$ 
  - register states  $\{|0\rangle_R, |1\rangle_R, \dots, |N_T\rangle_R\}$  represent no tick, 1 tick,  $\dots$ , and  $N_T$  ticks having occurred
  - $\rho_{CR}(0) = \rho_C^0 \otimes |0\rangle\langle 0|_R$
  - $\rho_{CR}(t) = \sum_{k=0}^{N_T} \tilde{\rho}_C^{(k)}(t) \otimes |k\rangle\langle k|_R$
- dynamics are given by a family of CPTP maps  $\mathcal{M}_{CR \rightarrow CR}^t(\cdot) = e^{\mathcal{L}_{CR} t}(\cdot)$  parametrized by coordinate time  $t \geq 0$ 
  - Lindbladian  $\mathcal{L}_{CR}(\cdot) = -i[\bar{H}, (\cdot)] + \sum_j \bar{L}_j(\cdot)\bar{L}_j^\dagger - \frac{1}{2}\{\bar{L}_j^\dagger\bar{L}_j, (\cdot)\} + \bar{J}_j(\cdot)\bar{J}_j^\dagger - \frac{1}{2}\{\bar{J}_j^\dagger\bar{J}_j, (\cdot)\}$ , where  $\bar{H} = H \otimes \mathbb{1}_R$ ,  $\bar{L}_j = L_j \otimes \mathbb{1}_R$ ,  $\bar{J}_j = J_j \otimes O_R$  with  $O_R = |1\rangle\langle 0|_R + |2\rangle\langle 1|_R + \dots + |N_T\rangle\langle N_T - 1|_R$



## Example: Ladder ticking clock

Choosing some orthonormal basis  $\{|j\rangle_C\}_{j=1}^d$  of  $\mathcal{H}_C$ , the clock is given by

- $L_j = \begin{cases} |j+1\rangle\langle j|_C, & \text{for } j \neq d \\ 0, & \text{for } j = d \end{cases}$
- $J_j = \begin{cases} 0, & \text{for } j \neq d \\ |1\rangle\langle d|_C, & \text{for } j = d \end{cases}$
- $H = 0$  and  $\rho_C^0 = |1\rangle\langle 1|_C$



Most accurate classical clock achieving  $R_k = kR_1$  with  $R_1 = d$  via reset

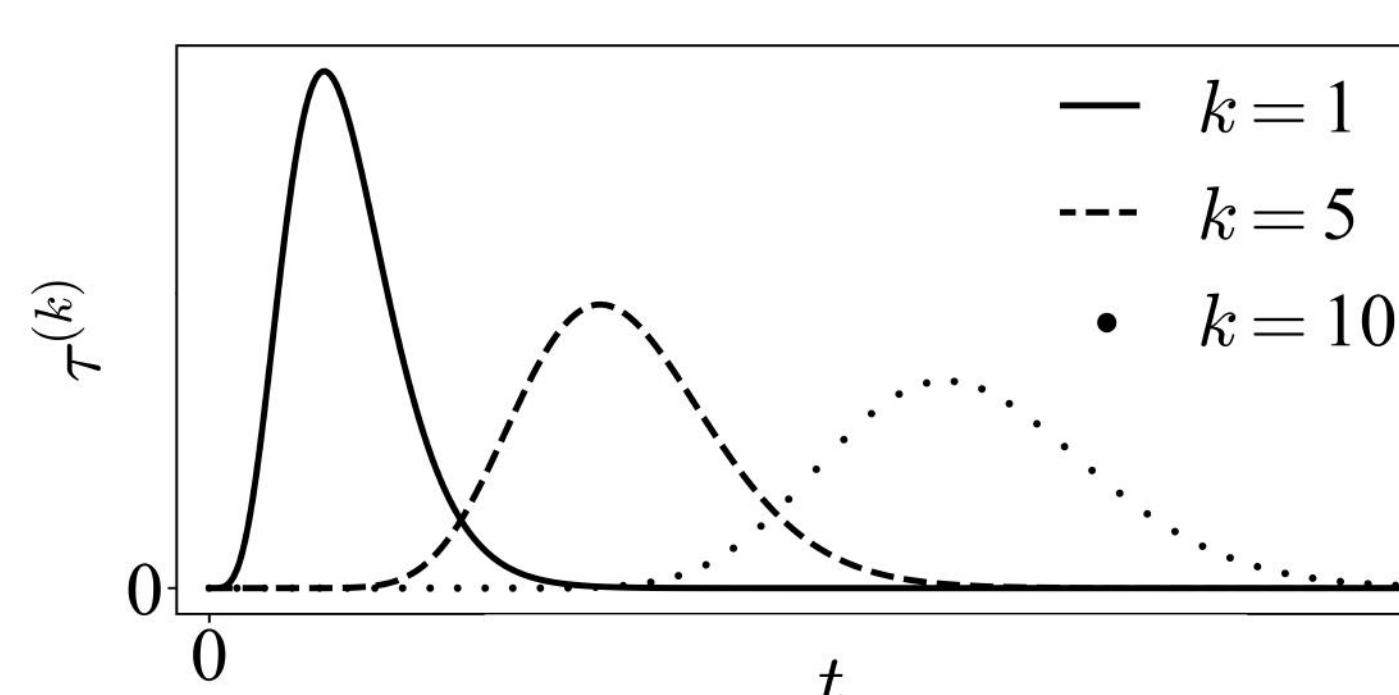
- classical  $\Leftrightarrow$  clockwork remains incoherent  $\rho_C(t) = \sum_j p_j(t)|j\rangle\langle j|_C$
- reset  $\Leftrightarrow \mathcal{C}_{\text{tick}}(\rho_C) = \rho_C^0 \forall \rho_C$ , where  $\mathcal{C}_{\text{tick}}(\cdot) = \sum_j J_j(\cdot)J_j^\dagger$

Quantum reset clocks can achieve a higher accuracy  $R_k = kR_1$  with  $R_1 \approx d^2$

## Measure of accuracy

based on delay functions  $\{\tau^{(k)}(t)\}_{k=1}^{N_T}$

- probability that  $(k-1)$  ticks occur during  $[0, t)$  and  $k$ th tick occurs during  $[t, t + \delta t]$  is  $\delta t \cdot \tau^{(k)}(t)$  with  $\delta t > 0$
- accuracy of  $k$ th tick is  $R_k = \mu_k^2 / \sigma_k^2$

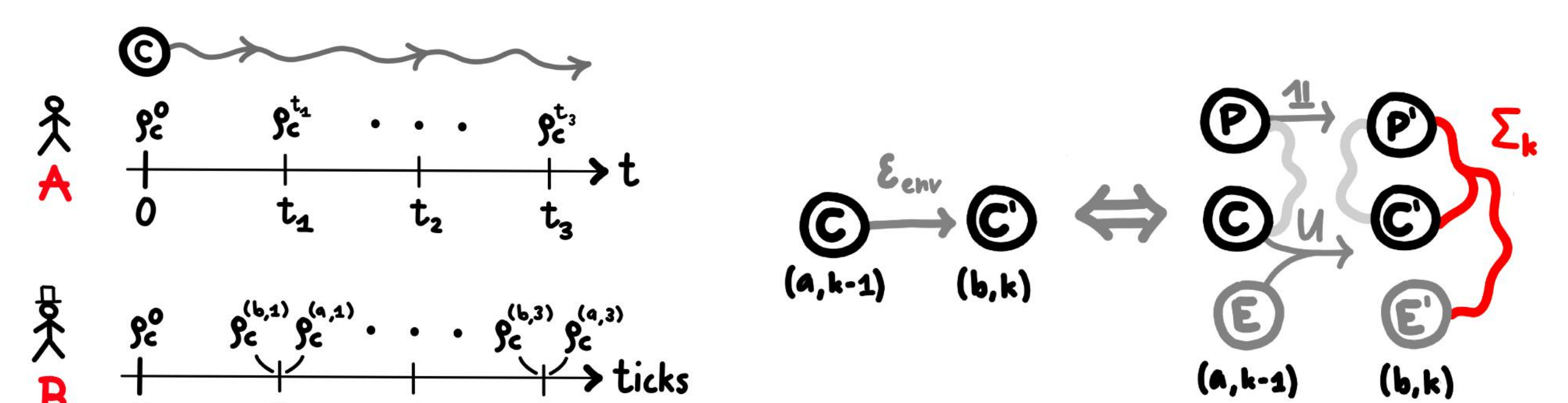


For ladder ticking clock:

## Measure of entropy production

$\Sigma_k = S(\rho_{CP}^{(b,k)})$  quantifies **exchanged information** between ticking clock and its environment during the  $k$ th tick

- entropy is **observer-dependent**  $\Rightarrow$  relevant observer (**B**) only has knowledge about ticks (not coordinate time  $t$  itself)



## Relationship between accuracy and entropy production

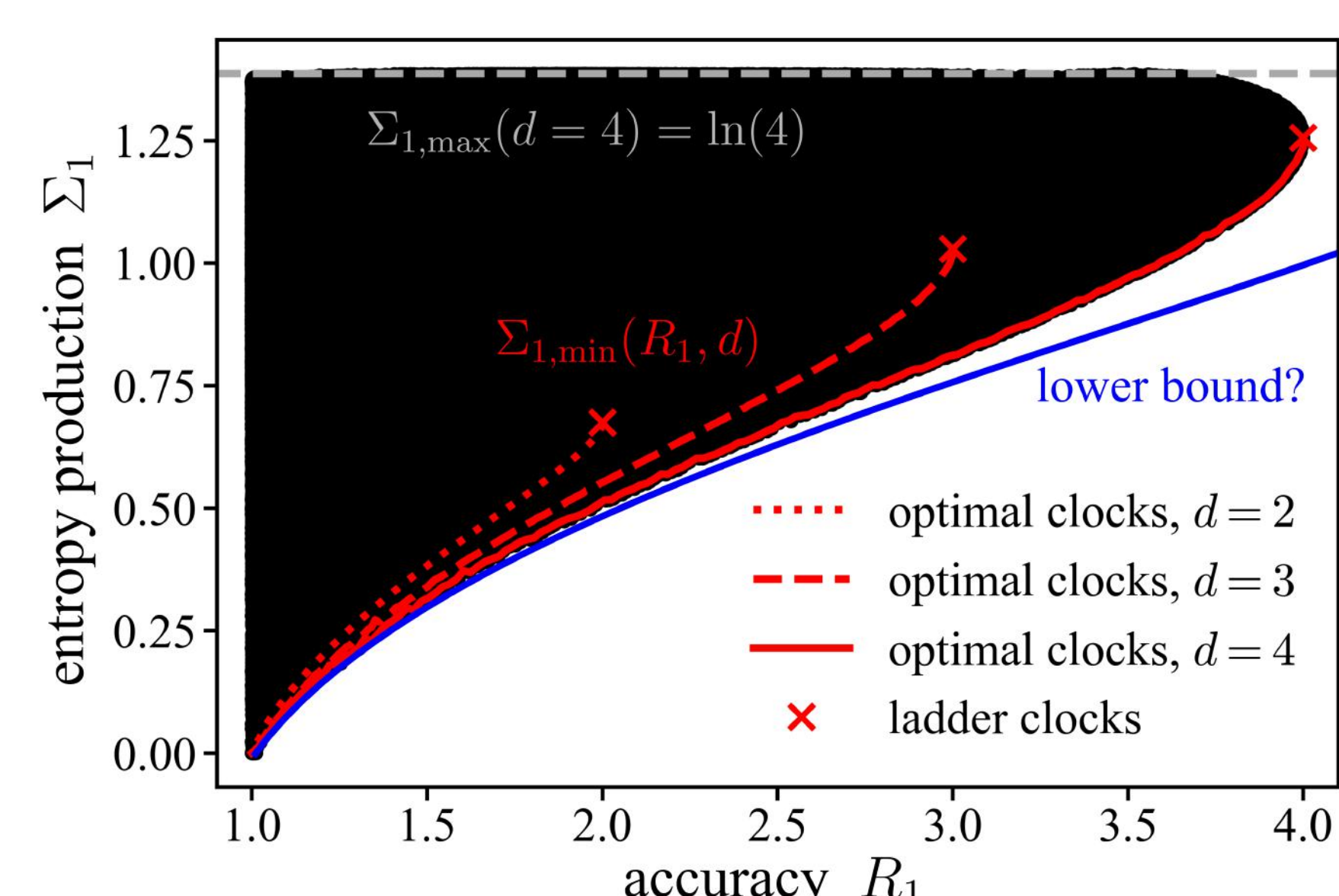
Vanishing entropy production per tick can be achieved by

- classical reset clock exhibiting Poissonian tick statistics with  $R_1 = 1$ 
  - $H = 0$ ,  $L_j = 0 \forall j$ ,  $J = \sqrt{\lambda}\mathbb{1}$ , and  $\rho_C^0 = |\Psi\rangle\langle\Psi|_C$
- "quasi-ideal" quantum reset clock with  $R_1 \approx d^2 \rightarrow \infty$ 
  - $H \neq 0$ ,  $L_j = 0 \forall j$ ,  $J_j \neq 0 \forall j$ , and  $\rho_C^0 = |\Psi\rangle\langle\Psi|_C$

$\Rightarrow$  proof crucially relies on vanishing no-tick generators  $\{L_j\}$

Note: Classical clocks with  $L_j = 0 \forall j$  achieve maximal accuracy of  $R_1 = 1$

Every classical ticking clock must produce a minimal amount of entropy per tick which increases with increasing accuracy



## References

- [P. Erker *et al.*, *Phy. Rev. X* **7**, 031022 (2017)]
- [M. P. Woods, *Quantum* **5**, 381 (2021)]
- [M. P. Woods *et al.*, *PRX Quantum* **3**, 010319 (2022)]

## Contacts

julian.arnold@unibas.ch  
mischa.woods@gmail.com

## Acknowledgments